

1) SPECIFICATION: 6+6 = 12% of the final mark (continuous assessment)

a) The following Haskell data declarations are useful for modelling natural and integers:

```
data Nat = Zero | Suc Nat
data Int = Z | S Int | P Int
deriving Show
```

Specify operation `addabs :: Int -> Nat -> Nat`, which returns the natural number representing the absolute value of the sum of an integer and a natural number. Although auxiliary operations are not permitted, to simplify the exercise we can assume that the `Int` number in the first parameter is ALWAYS IN NORMAL FORM. Examples:

```
addabs (P (P Z)) (Suc Zero) -> Suc (Suc Zero)
addabs (S (S Z)) (Suc Zero) -> Suc (Suc (Suc Zero))
```

b) Consider the following Haskell specification:

```
op23 :: (Ord a) => [a] -> [a]
op23 (x:y:s) = if (x>y) then s++[x] else x:(op23 (y:s))
op23 [] = []
```

Write the Haskell terms obtained after evaluating the following expressions:

```
op23 [1,2,3] ->
op23 [3,2,1] ->
op23 [1,3,2] ->
```

2) IMPLEMENTATION: 10+13 = 23% of the final mark (continuous assessment)

Complete the templates of the two methods below according to the following guidelines:

- The Java code must be EFFICIENT, ITERATIVE, using NO MORE VARIABLES than the ones declared in the templates and applying the `try(...)` catch (`TADVacioException ex`) (`ex.printStackTrace()`) command when needed;
- Auxiliary methods can not be used except the following ones declared in the `Lista<E>` interface for manipulating objects belonging to all classes implementing such interface:


```
public boolean EsVacia(); // checks if a list is empty
public void Añade(E e); // inserts an element in a list
public E Cabeza(); throws TADVacioException; // returns the last inserted element on a list
public Lista<E> Cola(); throws TADVacioException; // returns a list without its last inserted element
```
- Objects declared of type `LD<E>` can also access the attributes `private int longitud` (number of element of the list) and `private Nodo<E> NodoCabeza` (link the last inserted node in the list), as well as the constructor of the `Nodo<E>` class `public Nodo(E e, Nodo<E> n) {Info=e; Siguiente=n;}`.

a) public `LD(Lista<E> l1, LD<E> l2)`
{ `Nodo<E> aux=null;` }

..... // COMPLETE THE JAVA CODE HERE
This constructor of the `LD<E>` class generates a list with the elements of two lists `l1` and `l2` after reversing the order of the first one (i.e., `l1`). For instance, if `l1=[1,2,3]` and `l2=[4,5,6]` then, the content of object this must be `[3,2,1,4,5,6]`.

```
b) public static <E extends Comparable> Lista<E> op23(Lista<E> l1)
{ E x,y; LD<E> aux1,aux2; aux1=new LD<E>(); aux2=new LD<E>();
  .... // COMPLETE THE JAVA CODE HERE
  return new LD<E>(aux1,aux2); }
```

This method implements the previously specified `op23` operation outside all classes implementing the `Lista<E>` interface.

SOLUTIONS

1) a) `addabs :: Int -> Nat -> Nat`

```
addabs Z = Zero
addabs (S x) = Suc (addabs x)
addabs (P x) = Suc (addabs x)
```

b) `op23 :: (Ord a) => [a] -> [a]`

```
op23 (x:y:s) = if (x>y) then s++[x] else x:(op23 (y:s))
op23 [] = []
```

```
op23 [1,2,3] -> [1,2,3]
op23 [3,2,1] -> [1,3]
op23 [1,3,2] -> [1,3]
```

2) a) `public LD(Lista<E> l1, LD<E> l2) {`

```
    longitud = l2.longitud; NodoCabeza = l2.NodoCabeza;
    try { while (!l1.EsVacia()) {
        NodoCabeza = new Nodo(l1.Cabeza(), NodoCabeza);
        l1 = l1.Cola(); longitud++;
    } } catch (TADVacioException ex) { ex.printStackTrace(); }
```

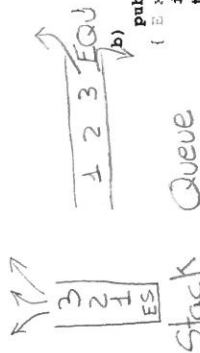
b) `public static <E extends Comparable> Lista<E> op23 (Lista<E> l1`

```
{ E x,y; LD<E> aux1,aux2; aux1=new LD<E>(); aux2=new LD<E>();
  if (!l1.EsVacia() || !l2.Cola().EsVacia()) return l1;
  try {
      x = l1.Cabeza(); l = l1.Cola(); y = l2.Cabeza();
      while (!l.Cola().EsVacia() && x.compareTo(y) <= 0) {
          aux1.Añade(x); x = y; y = l.Cabeza(); l = l.Cola();
      }
      if (x.compareTo(l.Cabeza()) > 0) { aux2.Añade(x); l = l.Cola();
          aux1.Añade(l.Cabeza()); l = l.Cola();
      }
      } catch (TADVacioException e) { return new LD<E>(); }
  return new LD<E>(aux1,aux2); }
```

JAVA CODE FOR TESTING:

```
LD<Integer> l1=new LD<Integer>(0);l1.Añade(3);l1.Añade(2);l1.Añade(1);
LD<Integer> l2=new LD<Integer>(0);l2.Añade(6);l2.Añade(5);l2.Añade(4);
System.out.println(" l1="+ l1 + " l2="+ l2 + " new constructor=" +
    new LD(l1,l2));

LD<Integer> l3=new LD<Integer>(0);l3.Añade(3);l3.Añade(2);l3.Añade(1);
System.out.println(" l3="+ l3 + " op23=" +op23(l1,l2));
l3=new LD<Integer>(0);l3.Añade(1);l3.Añade(2);l3.Añade(3);
System.out.println(" l3="+ l3 + " op23=" +op23(l1,l2));
l3=new LD<Integer>(0);l3.Añade(2);l3.Añade(3);l3.Añade(1);
System.out.println(" l3="+ l3 + " op23=" +op23(l1,l2));
```



- 1) Consider the following Haskell declarations useful for specifying a three-valued variant of Booleans based on just one NON FREE generator:

```
data Nat = Zero | Suc Nat
data BoolThreeNat = B3N Nat Nat
```

The infinite set of terms built with the **B3N** generator can be partitioned in three equivalence classes representing the following notions of truth:

- **True**: when the first parameter of the term is a **Nat** number strictly bigger than the second one. The canonical representative of this class is **B3N (Suc Zero) Zero**.
- **Maybe**: when the first parameter of the term is a **Nat** number coincides with the second one. The canonical representative of this class is **B3N Zero Zero**.
- **False**: when the first parameter of the term is a **Nat** number strictly smaller than the second one. The canonical representative of this class is **B3N Zero (Suc Zero)**.

Give the set of equations modelling these operations:

- a) **nfb3N :: BoolThreeNat -> BoolThreeNat**, which computes the normal form of a **BoolThreeNat** term without using auxiliary operations.
b) **andB3N :: BoolThreeNat -> BoolThreeNat**, which returns the conjunction of two **BoolThreeNat** terms. Calls to **nfb3N** as well as to an auxiliary operation (that must be also specified) are allowed.

Examples:

```
nfb3N (B3N (Suc (Suc Zero)) Zero) --> B3N (Suc Zero) Zero
nfb3N (B3N (Suc (Suc Zero)) (Suc (Suc (Suc Zero)))) --> B3N Zero (Suc Zero)
nfb3N (B3N (Suc (Suc Zero)) (Suc (Suc Zero))) --> B3N Zero Zero
andB3N (B3N (Suc (Suc Zero)) (B3N (Suc Zero) Zero)) --> B3N Zero Zero
andB3N (B3N Zero (Suc (Suc Zero))) (B3N (Suc Zero) Zero) --> B3N Zero (Suc Zero)
```

- 2) Consider the following Haskell specification:

```
queuexam :: Queue Int -> Queue Int
queuexam (q::x:y) = (q::y:x)
queuexam _ = queuexam (Eq::1:2)
```

- a) Write the Haskell terms obtained after evaluating the following expressions:

```
queuexam (Eq::3) --> Eq::3
queuexam (Eq::3::2::1) --> Eq::3,1:2
```

Implement this operation as two different Java methods with the following shapes:

- b) **HIGH LEVEL**: `public static Cola<Integer> queuexam (Cola<Integer> q)`
c) **LOW LEVEL**: `public void queuexam ()`

Use only the methods/attributes below according to their **HIGH/LOW** labels for each exercise.

```
public interface Cola<E>{
    public boolean EsVacia();
    public void EnCola(E x);
    public E Cabeza() throws TADVacioException;
    public Cola<E> Resto() throws TADVacioException;
}

public class Cola_Dinamica<E> implements Cola<E>{
    public Node<E> NodoCabeza; public Node<E> NodoFinal; -- LOW
    public Cola_Dinamica() {NodoCabeza=null;NodoFinal=null;} -- HIGH
}

public class Node<E> {
    public E Info; public Node<E> Siguiente;
    public Node(E e, Node<E> n){Info=e;Siguiente=n;}
}
```

SOLUTIONS

- 1) `data BoolThreeNat = B3N Nat Nat deriving Show`

```
a) nfb3N :: BoolThreeNat -> BoolThreeNat
nfb3N (B3N Zero Zero) = B3N Zero Zero
nfb3N (B3N (Suc Zero) Zero) = B3N (Suc Zero) Zero
nfb3N (B3N Zero (Suc Zero)) = B3N Zero (Suc Zero)
nfb3N (B3N (Suc x) (Suc y)) = nfb3N (B3N x y)
```

```
b) andB3N :: BoolThreeNat -> BoolThreeNat
andB3N x y = aux_andB3N (nfb3N x) (nfb3N y)
```

```
aux_andB3N :: BoolThreeNat -> BoolThreeNat -> BoolThreeNat
aux_andB3N (B3N (Suc Zero) Zero) x = x
aux_andB3N (B3N Zero (Suc Zero)) _ = B3N Zero (Suc Zero)
aux_andB3N (B3N Zero Zero) _ = B3N Zero Zero
aux_andB3N _ _ = B3N Zero Zero
```

- 2) `queuexam :: Queue Int -> Queue Int`
`queuexam (q::x:y) = (q::y:x)`
`queuexam _ = queuexam (Eq::1:2)`

```
a) queuexam (Eq::3) --> (Eq::2::1)
queuexam (Eq::3::2::1) --> (Eq::3::1::2)
```

```
b) public static Cola<Integer> queuexam (Cola<Integer> q)
{ Cola_Dinamica<Integer> sol=new Cola_Dinamica<Integer>();
  try {
    if (q.EsVacia() || q.Resto().EsVacia())
      { sol.EnCola(2);sol.EnCola(1); }
    else
      { while (!q.Resto().Resto().EsVacia())
          { sol.EnCola(q.Cabeza()); q=q.Resto(); }
        Integer x; x= q.Cabeza(); q=q.Resto();
        sol.EnCola(q.Cabeza());sol.EnCola(x);
      }
  } catch (TADVacioException ex) { System.out.println("TAD lineal vacio"); }
  return sol;
}
```

OTHER VERSION (which perhaps is more elegant)

```
public static Cola<Integer> queuexam (Cola<Integer> q)
{ Cola_Dinamica<Integer> sol=new Cola_Dinamica<Integer>();
  try { while (!q.Resto().Resto().EsVacia())
      { sol.EnCola(q.Cabeza()); q=q.Resto(); }
    Integer x; x= q.Cabeza(); q=q.Resto();
    sol.EnCola(q.Cabeza()); q=q.Resto();
    sol.EnCola(x);sol.EnCola(x); return sol;
  } catch (TADVacioException ex)
  { sol.EnCola(2);sol.EnCola(1); return sol;
  }
}
```

```
c) public void queuexam ()
{ if (NodoCabeza==null || NodoCabeza.Siguiente==null)
  { NodoFinal=new Node(1, null); NodoCabeza=new Node(2, NodoFinal); }
  else
  { Node<E> aux=NodoCabeza;
    while (aux.Siguiente!=NodoFinal) aux=aux.Siguiente;
    E x; x=aux.Info; aux.Info=NodoFinal.Info;NodoFinal.Info=x;
  }
}
```


$op22(q: x) s = x:/(op22 q s)$
 $op22 q (x:/y:/s) = y:/ (op22 q(x:/s))$
 $op22 _ _ = ES$

En Haskell NO desapila

$op22 (Equ: 1:2:3) (4: ES)$
 $3:/ (op22 Equ: 1:2) (4: ES)^{(1)}$
 $3:/2:/ (op22(Equ: 1) (4: ES))^{(2)}$
 $3:/2:/1:/ (op22(Equ) (4: ES))^{(2)}$
 $3:/2:/1:/ES$

$op22 Equ (4:15:/6:/ES)^{(2)}$
 $5:/ op22 Equ (4:/6:/ES)^{(2)}$
 $5:/6:/ op22 Equ (4:/ES)^{(2)}$
 $5:/6:/ES$

$op22 (Equ: 1:2:3) (4:/5:/6:/ES)$
 $3:/2:/1:/ op22 Equ (4:/5:/6:/ES)^{(2)}$
 $3:/2:/1:/5:/ op22 Equ (4:/6:/ES)^{(2)}$
 $3:/2:/1:/5:/6:/ op22 Equ (4:/ES)^{(2)}$
 $3:/2:/1:/5:/6:/ES$

a) High level in JAVA

```

public static <E> Pila<E> op22 (Cola<E> q, Pila<E> s) {
    Pila<E> s1; Cola<E> q1;
    Pila<E> s1 = Pila<E> s.clone(); { Para no modificar originales
    Cola<E> q1 = Cola<E> q.clone();
    try { if (!s1.EsVacia) s1 = s1.Desapila(); } → Si hay algo NO se puede acceder directamente
        while (!q1.EsVacia()) {
            s1.Apila(q1.Cabeza()) → Introduce el primer elemento de la cola
            q1 = q1.Resto() → q1 se actualiza con el resto de la cola
        } catch (...) {} ... {
    return s1;
    }
}

```

NodoCabeza → Pila
 q.NodoCabeza → Cola

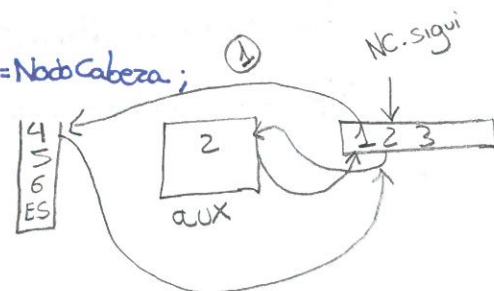
b) Low level in JAVA

```

public void op22 (CD<E> q) {
    Nodo<E> aux = null;
    if (NodoCabeza != null) NodoCabeza = NodoCabeza.Siguiente
    while (q.NodoCabeza != null) {
        aux = q.NodoCabeza.Siguiente; q.NodoCabeza.Siguiente = NodoCabeza;
        NodoCabeza = q.NodoCabeza; q.NodoCabeza = aux;
    }
}

```

"desapila el 4"
 4
5
6
ES



$subI \ Z \ (P \ Z) \rightarrow S(subI \ x \ y)$
 $subI \ Z \ (S \ Z) \rightarrow P(subI \ x \ y)$
 $subI \ (S \ Z) \ Z \rightarrow x$
 $subI \ (P \ Z) \ Z \rightarrow x$
 $subI \ (S \ Z) \ (S \ Z) \rightarrow subI \ x \ y$
 $subI \ (P \ Z) \ (P \ Z) \rightarrow subI \ x \ y$

{ Va quitando "S" y "P" hasta que se desigualen

$\text{mix21} :: \text{stack } a \rightarrow \text{Queue } a \rightarrow \text{Queue } a$
 $\text{mix21}(x:s) (q:y) = (\text{mix21}(y:s)q) :: x$
 $\text{mix21}(x:s) q = \text{mix } s \ q$
 $\text{mix21 ES } q = q$

$(q:y) \rightarrow$ último elemento
 $\hookrightarrow$ el resto

1
2

a) $\text{mix21}(1:2:ES) (EQ:3:4:5)$

- 1- No modifico la cola
- 2- Todo de la clase E

(1) $\text{mix21}(5:2:ES) (EQ:3:4) :: 1$

(1) $\text{mix21}(4:2:ES) (EQ:3:5) :: 1$

(2) $\text{mix21}(3:2:ES) EQ$ 4:5:1

(2) $\text{mix21}(2:ES) EQ$ $\rightarrow EQ:4:5:1$

(3) mix21 ES EQ

b) High level in JAVA

$\text{public static } \langle E \rangle \text{ Cola } \langle E \rangle \text{ mix21}(\text{Pila } \langle E \rangle s, \text{ Cola } \langle E \rangle q)$
 recibe 2 argumentos, stack y Cola

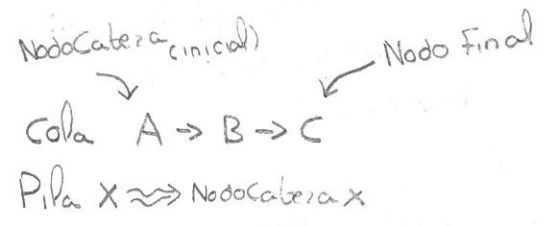
$\text{cola } \langle E \rangle q2 = (\text{Cola } \langle E \rangle) (q.\text{clone}());$ $\rightarrow$ No queremos tocar la cola original, creamos una nueva

$\text{if } (!s.\text{EsVacia}() \ \&\& \ !q2.\text{EsVacia}()) \}$ $\rightarrow$ Comprueba que no haya elementos vacíos

$q2 = q2.\text{Resto}();$ $\rightarrow$ Toma todos menos la cabeza "EQ:3:4:5"

$q2.\text{EnCola}(s.\text{Tope}());$ $\rightarrow$ Toma lo de arriba del stack y lo añade a la queue

$\{$
 $\text{return } q2;$
 $\{$



c) Low level in JAVA instancia clase PD <E>

$\text{public void mix21}(\text{PD } \langle E \rangle s)$

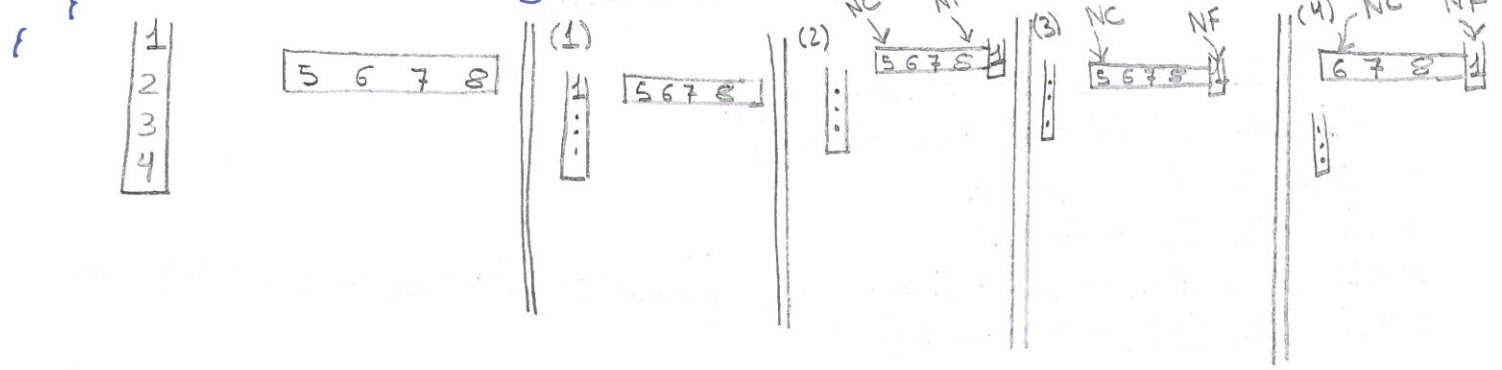
$\text{if } (s.\text{NodoCabeza} != \text{null} \ \&\& \ \text{NodoCabeza} != \text{null}) \}$ $\rightarrow$ Comprueba no nulo de stack y cola

(1) $s.\text{NodoCabeza}.\text{Sigiente} = \text{null};$ $\rightarrow$ Desconecta la cabeza del stack

(2) $\text{NodoFinal}.\text{Sigiente} = s.\text{NodoCabeza};$ $\rightarrow$ Actualiza el puntero final a cola tras añadir la antigua cabeza del stack

(3) $\text{NodoFinal} = \text{NodoFinal}.\text{Sigiente};$ $\rightarrow$ Mete 1º elemento de stack

(4) $\text{NodoCabeza} = \text{NodoCabeza}.\text{Sigiente};$



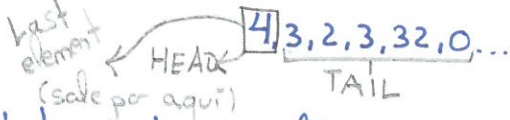
Fundamental Abstract Data Types

List ADT:

Linear sequence of elements with order relation, with a set of operations to manipulate.

Head: 1st element in the list, last introduced. LIFO (Last In First out)

Tail: The rest elements



data List $a = \text{Nil} \mid \text{Cons } a (\text{List } a)$

List $a \rightarrow [a]$

Nil $\rightarrow []$ Empty list

Cons $\rightarrow :$ (infix) Free generator (conservative)

- Operations no generators

Empty :: $[a] \rightarrow \text{Bool}$ // Checks if true

Head :: $[a] \rightarrow a$ // Returns head

Tail :: $[a] \rightarrow [a]$ // Returns tail

$\begin{cases} \text{head } (h:_) = h \\ \text{tail } (_:t) = t \end{cases}$

Length :: $[a] \rightarrow \text{int}$

Concatenate :: $[a] \rightarrow [a] \rightarrow [a]$ // Seen as ++

$\begin{cases} \text{length } (h:t) = 1 + \text{length } t \\ \text{concatenate } (h:t) s = h: \text{concatenate } t s \\ \text{reverse } (h:t) = \text{concatenate } (\text{reverse } t) [h] \end{cases}$

Reverse :: $[a] \rightarrow [a]$

Last :: $[a] \rightarrow a$

Elem :: $(\text{Eq } a) \Rightarrow a \rightarrow [a] \rightarrow \text{Bool}$

Extract :: $(\text{Eq } a) \Rightarrow a \rightarrow [a] \rightarrow [a]$

$\begin{cases} \text{elem } x (h:t) = (x == h) \parallel \text{elem } x t \\ \text{extract } x (h:t) = \text{if } (x == h) \text{ then extract } x t \text{ else } h: \text{extract } x t \end{cases}$

Take :: $\text{Int} \rightarrow [a] \rightarrow [a]$ // Only "saves" the Int elements from left

Cut :: $\text{Int} \rightarrow [a] \rightarrow [a]$ // Deletes Int elements from left

$\begin{cases} \text{take } _ [] = [] \\ \text{take } n (h:t) = h: \text{take } (n-1) t \\ \text{cut } n (h:t) = \text{cut } (n-1) t \end{cases}$

Substitute :: $(\text{Eq } a) \Rightarrow a \rightarrow a \rightarrow [a] \rightarrow [a]$

// substitute x y: changes all X's in [a] by Y's

$\begin{cases} \text{substitute } x y (h:t) = \text{if } h == x \\ \text{then } y: \text{substitute } x y t \\ \text{else } h: \text{substitute } x y t \end{cases}$

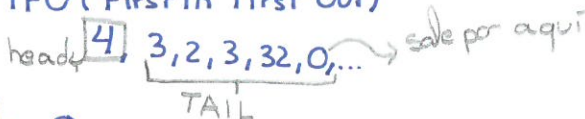
Sum :: $[\text{Int}] \rightarrow \text{Int}$ // Sum of all numbers

Greater :: $[\text{Int}] \rightarrow \text{Int}$ // greater :: $(\text{Ord } a) \Rightarrow [a] \rightarrow a \dots$

// greater (x:y:t) = if x > y then grt(x:t), else grt(y:t)

Queue ADT

FIFO (First In First out)



data Queue $a = \text{EQ} \mid \text{Put } (\text{Queue } a) a$

EmptyQ :: $\text{Queue } a \rightarrow \text{Bool}$

// EmptyQ EQ = True

// EmptyQ _ = False

First :: $\text{Queue } a \rightarrow a$

// First (EQ::x) = x

// First (c::x) = First c

Recordar que en color el primero está al final

$\text{first } (\text{EQ}::3::5::2::7) \rightarrow 3$

$\text{rest } (\text{EQ}::3::5::2::7) \rightarrow (\text{EQ}::5::2::7)$

$\text{Put } (\text{Put } (\text{Put } (\text{Put } (\text{EQ} 2) 7) 5) 7) \rightarrow \text{EQ}::2::7::5::7$

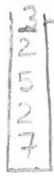
Rest :: $\text{Queue } a \rightarrow \text{Queue } a$

// rest (EQ::x) = EQ

// rest (c::x) = (rest c)::x

Stack ADT (Pilas)

Based on sequential access, represented vertically



TOP (first element)

A "pop" is necessary to access the remaining elements

$3 : (2 : (5 : \dots : (ES))) \rightarrow 3 : 2 : \dots : ES$

Empty S :: Stack a $\rightarrow$ Bool Top :: Stack a $\rightarrow$ a Pop :: Stack a $\rightarrow$ Stack a
 $\hookrightarrow \simeq$ Head $\hookrightarrow \simeq$ Tail

$\rightarrow$ Se debe desapilar antes de acceder

elimina :: a $\rightarrow$ Stack a $\rightarrow$ Stack a

elimina x ES = ES

elimina x (y : p) = if x == y then p, else y : (elimina x p)

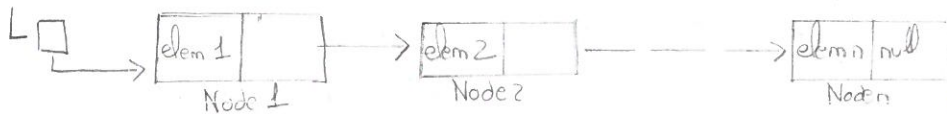
Toca el interior de la pila
 $\hookrightarrow +$ apila x p = x : p
 $\hookrightarrow$ elimina x (y : p) = if x == y then p else apila y (elimina x p)

Implementation of lists

In java, lists are interface List<E>, E is the class of elements in the list.

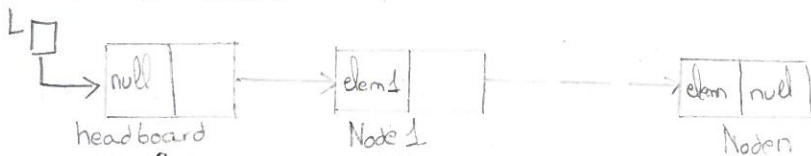
• Dynamic implementation

$\hookrightarrow$ must be extended for any class implementation list



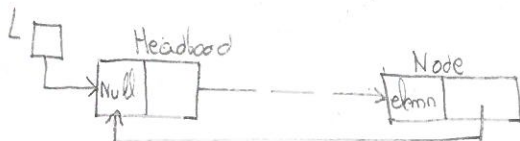
Node 1, ..., Node N instances of Node<E> class,
 L instance of List<E> interface.
 Elem 1, elem N, ... elements of class E to be stored in the list

• With headboard

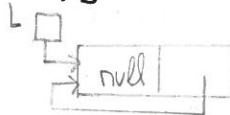


Avoids asymmetry in first node, 1st unique without predecessor

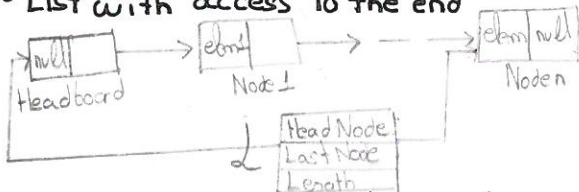
• Circular lists



• Empty list



• List with access to the end



Useful for concatenations or inserting after the last element avoiding to go through all nodes

Dynamic implementation of Queues

2324 First Progress Test

data Nat = Zero | Suc Nat $\rightarrow$ Números Naturales $\Rightarrow$ (Num ≥ 0)

data Int = Z | S Int | P Int $\rightarrow$ Números enteros (Positivos y negativos)

addabs :: Int $\rightarrow$ Nat $\rightarrow$ Nat $\rightarrow$ Metemos un entero y un natural, obteniendo otro natural

① addabs Z x = x $\rightarrow$ Siempre será positivo, ya que el primer valor (Inte) podría tomar cualquier valor, pero está puesto a cero (Z) y el segundo (Nat) sobre es ($x \geq 0$)

Toy example:

addabs Z (Suc Zero) $\rightarrow$ Suc Zero

② addabs (S Z) x = Suc (addabs x y) $\rightarrow$ Siempre será positivo por el mismo motivo.

Toy example

La recursividad "despeja" una S, luego Z = x; x = y

addabs (S (S Z)) (Suc Zero) $\rightarrow$ (Suc (Suc (Suc Zero)))

③ addabs (P Z) (Suc Zero) = addabs x y

addabs (P Z) (Zero) = Suc (addabs x Zero) $\rightarrow$ Every time the code finds a P, it will subtract 1 ordado!

addabs :: Int $\rightarrow$ Nat $\rightarrow$ Nat

addabs Z x = x

addabs (S Z) y = Suc (addabs x y)

addabs (P Z) (Suc Zero) = addabs x y

addabs (P Z) Zero = Suc (addabs x y)

$$-2 + 0 = 1 + (-1 \ 0)$$

$$-1 + 0 = 2 \ (0 \ 0) = \text{Suc}(\text{Suc Zero})$$

op23 :: (Ord a) $\rightarrow$ [a] $\rightarrow$ [a]

op23 (x:y:z) = if (x > y) then s++[x] else x:(op23 (y:s))

op23 [] = []

"s" concatena elemento "x"

op23 [1,2,3] $\rightarrow$ [1,2,3]

op23 [3,2,1] $\rightarrow$ [1,3] $\rightarrow$ 3 > 2; [1] ++ [3]; [1,3]

op23 [1,3,2] $\rightarrow$ [1,3] $\rightarrow$ 1 > 1; 3 > 2; [1] ++ [3]; [1,3]

$\hookrightarrow$ 1 : (op23 [3:2:[1]]) $\Rightarrow$ 1 : [] ++ [3]

L1 y L2, dif class

Using times :: Nat $\rightarrow$ Nat $\rightarrow$ Nat // times a $\rightarrow$ a $\rightarrow$ a from practice 1

times (done)

times : a $\rightarrow$ a $\rightarrow$ a ## We use a parameter, we obtain one (a)

facto : Nat $\rightarrow$ Nat ### We introduce one natural, we obtain another one

facto : Zero $\rightarrow$ sucZero ### If we obtain zero, the result is one $\rightarrow 0! = 1$

facto : Suc (x) $\rightarrow$ times (fact x) (suc x) ### The natural is introduced

$\downarrow$
Num $\neq 0$

$\checkmark$
 $n! = n \cdot (n-1)$ recursively

by example:

facto Suc (suc x) $\rightarrow$ times (fact (Suc x)) Suc (suc x)

$2! \rightarrow 1 \cdot 2$

If higher than 2, this result would be save. For example 3 would be $6 \cdot 1$ at the end, after $3 \cdot 2$

2021) 6) public static <E> Cola<E> mix21 (Pila<E> s, Cola<E> q)

```

    Cola<E> q1 = (Cola<E>) q1.clone();
    if (!s.EsVacia() || !q1.EsVacia())
        q1 = q1.Resto();
    q1.EnCola(s.Tope());

```

```

    {
        return q1;
    }

```

- 1 - Clono
- 2 - Hay vacios?
- 3 - Apilo / desapilo lo nuevo
- 4 - Devuelvo

2022) a) public static <E> Pila<E> op22 (Cola<E> q, Pila<E> s)

```

    Pila<E> s1 = (Pila<E>) s.clone();
    Cola<E> q1 = (Cola<E>) q.clone();
    if (!s1.EsVacia()) s1 = s1.Desapila();
    while (!q1.EsVacia())
        s1.Apila(q1.Cabeza());
    q1 = q1.Resto();

```

```

    {
        return s1;
    }

```


11

1

1) SPECIFICATION: $10+3 = 13\%$ of the final mark (continuous assessment)

a) The following data declaration (based on NON FREE generators) is useful for modelling integers: $\text{data Inte} = \mathbb{Z} \mid \text{S Inte} \mid \text{P Inte}$ deriving Show. WITHOUT using auxiliary operations, specify a subtraction operation $\text{subI} :: \text{Inte} \rightarrow \text{Inte} \rightarrow \text{Inte}$, such that: if the inputs are not in normal form, then the output doesn't need to be in normal form but, if the inputs are in normal form, then the output must be also in normal form. Examples:

$\text{subI (P (P (S Z))) (S (P (S Z))))} \rightarrow \text{P (P (P (P (S Z))))}$
 $\text{subI (P Z) (S (S Z)))} \rightarrow \text{P (P (P Z))}$

b) Consider the following Haskell specification:

```
infix 5 :: data Stack a = ES | a :: (Stack a) deriving Show
infix 1 5 :: data Queue a = EQ | (Queue a) :: a deriving Show
op22 :: Queue a -> Stack a -> Stack a
```

$\text{op22 (q::x) s} = \text{x} :: (\text{op22 q s})$
 $\text{op22 q} = \text{y} :: (\text{op22 q (x::a)})$
 $\text{op22} = \text{ES}$

Write the Haskell terms obtained after evaluating the following expressions:

$\text{op22 (SQu::1..2..3) (4::ES)} \rightarrow 3::1/2/1/4::ES$
 $\text{op22 SQu (4::5::6::ES)} \rightarrow 5::1/6::ES$
 $\text{op22 (SQu::1..2..3) (4::5::6::ES)} \rightarrow 3::1/2/1/4/5::1/6::ES$

2) IMPLEMENTATION: $11+11 = 22\%$ of the final mark (continuous assessment)

Implement the previous op22 operation as two EFFICIENT Java methods according to the following guidelines: 1) no more variables than the ones declared in the templates are permitted; 2) the try/catch (TADVacioException ex) {ex.printStackTrace();} command must be unchanged after calling such method and 4) the LOW level method will be able to alter the content of the input parameter as well as the content of object this inside the PD class.

a) HIGH LEVEL (use only the methods declared in the Pila and Cola interfaces):

```
public static <E> Pila<E> op22 (Cola<E> q, Pila<E> s)
{ Pila<E> s1; Cola<E> q1; ... }
```

b) LOW LEVEL (use only the attributes of the CD, PD and Nodo classes):

```
public void op22 (CD<E> q)
{ Nodo<E> aux=null; ... }
```

```
public interface Pila<E>{
    public boolean EsVacia();
    public void Apila(E e);
    public E Tопо(); throws TADVacioException;
    public Pila<E> Desapila();throws TADVacioException;
    public Object clone();
}

public interface Cola<E>{
    public boolean EsVacia();
    public void EnCola(E x);
    public E Cabeza() throws TADVacioException;
    public Cola<E> Resto() throws TADVacioException;
    public Object clone();
}

public class PD<E> implements Pila<E>{
    public class CD<E> implements Cola<E>{
        public class Nodo<E> NodoCabeza; public class Nodo<E> NodoFinal;
        public class Nodo<E> {
            public E info;
            public Nodo<E> siguiente;
        }
    }
}
```

SOLUTIONS

a) $\text{subI x z} = \text{x}$
 $\text{subI (S x) (S y)} = \text{subI x y}$
 $\text{subI x (S y)} = \text{P (subI x y)}$
 $\text{subI (P x) (P y)} = \text{subI x y}$
 $\text{subI x (P y)} = \text{S (subI x y)}$

b) $\text{op22 (Eq::1..2..3) (4::ES)} \rightarrow 3::1/2::1/4::ES$
 $\text{op22 EQ (4::5::6::ES)} \rightarrow 5::1/6::ES$
 $\text{op22 (Eq::1..2..3) (4::5::6::ES)} \rightarrow 3::1/2::1/4/5::1/6::ES$

```
public static <E> Pila<E> op22 (Cola<E> q, Pila<E> s)
try {
    Pila<E> s1 = (Pila<E>)s.clone(); Cola<E> q1 = (Cola<E>)q.clone();
    if (!s1.EsVacia()) s1 = s1.Desapila();
    while (!q1.EsVacia()) {
        s1.Apila(q1.Cabeza());
        q1 = q1.Resto();
    }
    catch (TADVacioException ex) { ex.printStackTrace(); }
    return s1;
}
```

```
b)
public void op22 (CD<E> q)
{ Nodo<E> aux=null;
  if (NodoCabeza!=null) NodoCabeza=NodoCabeza.siguiente;
  while (q.NodoCabeza!=null)
      {aux=q.NodoCabeza.siguiente;q.NodoCabeza.siguiente=NodoCabeza;
       NodoCabeza=q.NodoCabeza;q.NodoCabeza=aux;}
}
```

CD CODE FOR TESTING:

```
PD<Integer> q1=new CD<Integer>();
q1.EnCola(1); q1.EnCola(2);q1.EnCola(3);
PD<Integer> s1=new PD<Integer>();
s1.Apila(6); s1.Apila(5); s1.Apila(4);

System.out.println("HIGH op22 (Eq::1..2..3) (4::5::6::ES) --> 3 :: 2
:: 1 :: 5 :: 6 :: ES " + op22(q1,s1));

s1.op22(q1);
System.out.println("LOW op22 (Eq::1..2..3) (4::5::6::ES) --> 3 :: 2
:: 1 :: 5 :: 6 :: ES " + s1);
```



Si stack no vacío

```
public void op22 ( <D <E > q )
  Nodo <E> aux = null;
```

```
  if ( nodoCabeza != null ) {
    nodoCabeza = nodoCabeza . siguiente
    "Desapilamos el 4"
    "Mientras haya elementos en cola"
```



```
  while ( q . nodoCabeza != null ) {
    aux = q . nodoCabeza . siguiente ;
    q . nodoCabeza . siguiente = nodoCabeza ;
```

Se desconecta la cabeza de Queue (1)



```
  nodoCabeza = q . nodoCabeza ;
  "Se mete el 1 en la pila"
```

```
  q . nodoCabeza = aux ;
```



}

1) a) $\text{addabs} :: \text{Int} \rightarrow \text{Nat} \rightarrow \text{Nat}$

$$\text{addabs } z \ x = x$$

$$\text{addabs } (S \ z) \ y = \text{Suc}(\text{addabs } x \ y)$$

$$\text{addabs } (P \ z) \ (\text{Suc } y) = \text{addabs } x \ y$$

$$\text{addabs } (P \ z) \ y = \text{Suc}(\text{addabs } x \ y)$$

b) $\text{op23} :: (\text{Ord } a) \Rightarrow [a] \rightarrow [a]$

$\text{op23 } (x:y:s) = \text{if } (x > y) \text{ then } s ++ [x]$

$\text{op23 } [] = [] \text{ else } x : (\text{op23 } (y:s))$

$\text{op23 } [1,2,3]^{(1)}$

$1 : (\text{op23 } (2:3))^{(1)}$

$1 : 2 : (\text{op23 } (3:_))^{(1)}$

$[1:2:3]$

$\text{op23 } [3,2,1]$

$[3:1]^{(2)}$

$\text{op23 } [1,3,2]$

$1 : (\text{op23 } (3:2))^{(2)}$

$1 : 3$

$\hookrightarrow \text{NO hay } s$

2) $L1 = 1,2,3 \quad L2 = 4,5,6 \rightarrow [3,2,1,4,5,6]$ Necesito: 1- Nuevo nodo 2- Ampliar long $\hookrightarrow$ introducir L1 a L2

$\text{public LD } (L1 <E> L1, LD <E> L2)$

$\text{longitud} = L2.\text{longitud}; \quad \text{NodoCabeza} = L2.\text{NodoCabeza};$

$\text{try } \{ \text{while } (L1.\text{EsVacia}()) \}$

$\text{NodoCabeza} = \text{new Nodo}(L1.\text{Cabeza}(), \text{NodoCabeza});$

$\{ L1 = L1.\text{Cola}(); \text{longitud}++; \}$

$\{ \text{catch } (\dots) \} \dots \}$

$\hookrightarrow$ Impartentísimo!!! Para recorrer array y quitar usado

$\text{public static } <E> \text{Pila } <E> \text{Op22 } (Cola <E> q, Pila <E> s)$

$\text{Pila } <E> s1; \text{Cola } <E> q1;$

$s1 = (Pila <E> s).\text{clone}(); \quad q1 = (Cola <E> q).\text{clone}();$

$\text{try } \{ \text{if } (!s1.\text{EsVacia}()) \} s1 = s1.\text{Desapila}();$

$\text{while } (!q1.\text{EsVacia}())$

$s1.\text{Apila}(q1.\text{Cabeza});$

$q1 = q1.\text{Resto}();$

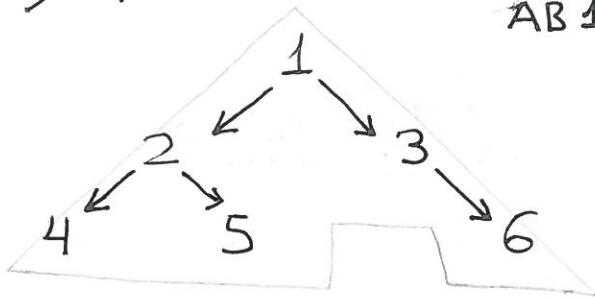
$\{ \text{catch } (\dots) \} \dots \{$

$\text{return } s1;$

$\}$

2/a)

AB 1 (AB 2 (AB 4 AV AV) (AB 5 AV AV) (AB 3 AV (AB 6 AV AV)



Output list: [6, 1, 0]

⇒ [6 | 1 | 0] → Y el 0???
↳ el 1??

40% BC → esVacio?
son así } check left recurs → Calling A i
check right, recurs → Calling A d { Tomar, el max tiene un hueco (l > r) || (r > l)

public static <E> Lista <Integer> hole (Arbol Binario <E> t)

int l, r; Lista <Integer> sol, left, right; sol = new LD() <> ();

if (t.EsVacio()) sol.Añade(0); // Indica que está vacío

else try { left = hole(t.SubArbol Izq()); right = hole(t.SubArbol Dcho());

// El método hole cuenta los nodos, los almacena en left/right
↳ // Los almacena en la cabeza

l = left.Cabeza(); r = right.Cabeza();

if (l < r) { sol = left.Cola(); sol.Añade(0); }

else if (l > r) { sol = right.Cola(); sol.Añade(1); }

sol.Añade(l+r+1)

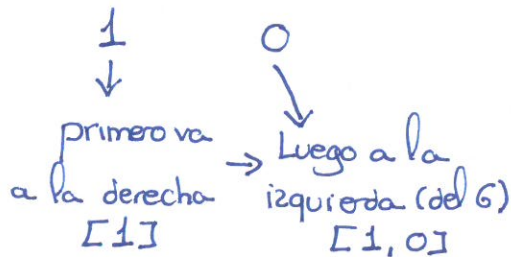
{ catch (TAD Vacio Excepcion ex) { ex.print... (); }

return sol;

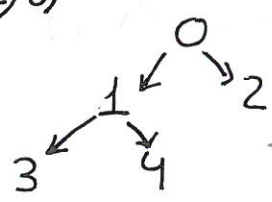
}

Condición previa en enunciado

6



2) b)



NotoGrafo <E>

Nodes

0	1	2	3	4			
---	---	---	---	---	--	--	--

false false

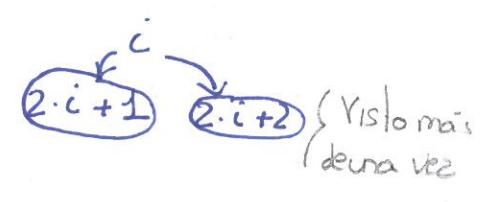
→ null
→ null...

Adyacentes

	0	1	2	3	4
0		T	T		
1	T			T	T
2	T				
3		T			
4		T			

21-22

No general trees



PN (12 % 5) * 4

↳ * (% 12 5) 4 = * % 12 5 4

*													
%		%		12		5	%	5	%	2		2	
12		12		5		4		4		4		4	
5		5		4		4		4		4		4	
4		4		4		4		4		4		4	
exp	aux	exp	aux	exp	aux	exp	aux	exp	aux	exp	aux	exp	aux

```

try { while (!exp. EsVacia())
    if (exp. Tope() instanceof Double && (!aux. EsVacia() && aux. tope() instanceof Double))
        x = (Double) exp. tope();
        exp = exp. Desapila();
        y = (Double) aux. tope();
        aux = aux. Desapila();
        op = (String) aux. tope();
        aux = aux. Desapila();
        exp. Apila (op Eval(x, op, y));
    else { aux. Apila (exp. tope());
        exp = exp. Desapila();
    }
    { catch (Tad) } ... {
    { return (Double) aux. tope();
    {

```

if (!t. EsVacio())

```

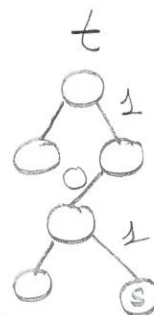
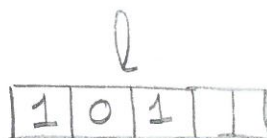
try { if (!l. EsVacio())
    if (t. subordDer(). EsVacio()) p = t. Raiz ();
    else if (l. Cabeza() == 0) {
        p = metodo (t. subordIzq(), l. Cola());
    } else { p = metodo (t. subordDer(), l. Cola());
    }
    { catch (...) } ... {

```

{ return p;

{ Pasos para crear un Huffman tree

- Paso lista y árbol
- Compruebo que el 2º elem. de la lista de árboles no sea nulo
- Lo es? Devuelvo el árbol único con sal = l. Cabeza(); 129 perc
- No lo es = Crea ABD con new ABD (raiz1 + raiz2, Arb1, Arb2)
- Inserto ABD en lista x = insert (ABD, x);



Java interfaces that can be used in this exam

```
public interface Lista<E> {
    public boolean EsVacia();
    public void Añade(E x);
    public E Cabeza() throws TADVacioException;
    public Lista<E> Cola() throws TADVacioException;
    public void Concatena(Lista<E> l);
    public Object clone();
}

public interface ArbolBinario<E> extends Arbol<E>
{
    public boolean EsVacio();
    public E Raiz() throws TADVacioException;
    public ArbolBinario<E> SubArbolIzqdo() throws TADVacioException;
    public ArbolBinario<E> SubArbolDcho() throws TADVacioException;
}
```

PN

$$2 + ((12 \% 5) * 3)$$

$$2 + ((12 \% 5) * 3)$$

$$+ 2 \left[(12 \% 5) * 3 \right]$$

$$+ 2 * \left[(12 \% 5) * 3 \right]$$

$$+ 2 * \% 12 5 3$$

RPN

$$2 + ((12 \% 5) * 3)$$

$$2 ((12 \% 5) * 3) +$$

$$2 \left[(12 \% 5) * 3 \right]$$

$$2 \left[(12 \% 5) * 3 \right] +$$

$$2 \left[(12 \% 5) * 3 \right] * +$$

$$2 12 5 \% 3 * +$$

1) [Reverse Polish Notation] In this exercise we are concerned with the second practice devoted to evaluate an arithmetic expression stored into a general stack once expressed in RPN (*reverse Polish notation*). Given the arithmetic expression $2 + ((12 \% 5) * 3)$, complete the following Java code for initializing a stack representing the same expression in reverse Polish notation:

```
Pila s = new PD(); s.APila( + ); s.APila( * ); s.APila( 3 );
s.APila( % ); s.APila( 5 ); s.APila( 12 ); s.APila( 2 );
```

In the table below, use only the cells required to represent the contents of stacks **exp** and **aux** after finishing each iteration when calling the method explained in class with the previous stack:

ITE. 1	ITE. 2	ITE. 3	ITE. 4	ITE. 5	ITE. 6	ITE. 7	ITE. 8	ITE. 9	ITE. 10
12									
5	5								
%	%	%							
3	3	3	5	3	3				
*	*	12	*	12	*	2	*	2	6
+	2	+	2	+	2	+	2	+	2
exp	aux	exp	aux	exp	aux	exp	aux	exp	aux

2) [Huffman trees] Complete in the next page the Java code of method **code01** which receives a Huffman tree **t** with at least two leaves and the ASCII code of a symbol **s** and returns a list only containing 0's and 1's representing the Huffman code of **s** if such symbol is included in the tree **t**. Otherwise, the output must be the list **[-1]**. The number of visited nodes in the tree must be minimized.

3) [Floyd's algorithm] A floydij in a weighted directed graph is a path connecting two NON ADJACENT nodes **i** and **j** and verifying that its weight is: 1) SMALLER than the weight of any other path between **i** and **j** and 2) BIGGER than the double of the weight of its first arrow. Complete in the next page the Java code of method **floydijs** which returns the number of floydijs found on a graph represented by its adjacency matrix.

```
public static Lista<Integer> code01 (ArbolBinario<Integer> t, Integer s)
{   Lista<Integer> l, laux; l=new LD(); ArbolBinario<Integer> aux=new ABD<>();
```

```
    return l;
}
```

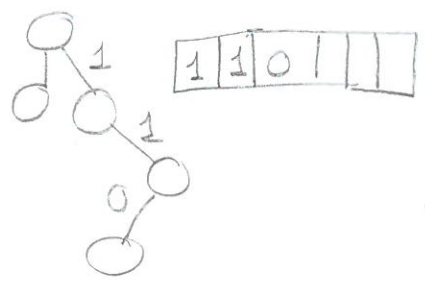
```
public static int floidijs (int[][] adj)
{int inf = Integer.MAX_VALUE, n = adj.length, i = 0, j = 0, k = 0, num = 0;
  int[][] paw = (int[][])adj.clone();   Lista<Integer>[][] pan = Floyd(paw);
```

```
    return num;
}
```

```

public static int percent(AB<int> t, lista<int> l)
int p = -1; AB<int> aux = new AB<>();
if (!t.EsVacio())
try { if (!l.EsVacia())
      if (t.subArbolDer.EsVacio()) p = t.Raiz();
    else if (l.Cabeza() == 0) p = percent(t.subArbolIzq(), l.Cola());
    else p = percent(t.subArbolDer(), l.Cola());
  }

```



```

public static ABD<int> huffmanTree(lista<ABD<int>> a)

```

```

try { while (!a.Cola().EsVacia())
      ABD<int> z = new ABD(a.Cabeza().Raiz() + a.Cola().Cabeza().Raiz(), a.Cabeza(), a.Cola().Cabeza());
      a = a.Cola().Cola();
      a = insert(ABD, a);
    } catch (...) {} ... {
  }

```

```

public static lista code01(AB<int> t, lista l)

```

```

try { if (t.subArbolDer().EsVacio())
      if (t.subArbolIzq().Raiz() != s) l.Añadir(-1);
    else l = code01(t.subArbolIzq(), l);
      if (l.EsVacia() || l.Cabeza() != -1) l.Añadir(0);
    else l = code01(t.subArbolDer(), s);
      if (l.EsVacia() || l.Cabeza() != -1) l.Añadir(1);
    }
  } catch (...) {} ... {

```

$$6 - (7 \times 2)$$

$$- 6(x + 2)$$

$$- 6 \times 72$$

```

try { while (!exp.EsVacio())

```

```

  if (exp.Tope() instanceof Double && (!aux.EsVacio() && aux.Tope() instanceof Double))
    x = exp.Tope(); y = aux.Tope(); op = aux.Tope(); exp.Apila(eval(x, op, y));
    exp = exp.Desapila(); aux = aux.Desapila(); aux = aux.Desapila();

```

```

  else aux.Apila(exp.Tope());

```

```

  exp = exp.Desapila(); return (Double) aux.Tope();

```

```

  { catch (...) {} ... {

```

```

  {

```


PN

	0%
	4
15	3
4	+
1	*
	+

RPN

$$351 - + 154\% * 1 +$$

$$\frac{15}{4}\%$$

4

3

3

3

7

7

21

21

22

22

3 exp 3 aux

$\{$

$7x((3-1)+(3-3))$

$x7((-3\ 1)+(-3\ 3))$

$x7(+(-3\ 1)(-3\ 3))$

$x7+-3\ 1-3\ 3$

s. Apila (3); s. Apila (3); s. Apila (-); ...

+
3
1
-
3
3 x
exp aux

+
3
1
-
3
3 x
exp aux

2
+
3
3 x
exp aux

-
3
1
-
3
3 x
exp aux

3
+
7
3 x
exp aux

-
3
1
-
3
3 x
exp aux

3
+
7
3 x
exp aux

1
3
+
7
3 x
exp aux

0
exp

1
3
+
7
3 x
exp aux

2
0
exp

2
+
7
3 x
exp aux

0
exp

2
+
7
3 x
exp aux

2
0
exp

7
3 x
exp aux

14
exp aux

```

try { while ( ! exp. EsVacia() ) {
    if ( exp. Tope() instanceof Double && ( ! aux. EsVacia() && aux. tope() instanceof Double ) ) {
        x = (Double) exp. Tope();
        exp = exp. Desapila();
        y = (Double) aux. Tope();
        aux = aux. Desapila();
        op = (String) aux. Tope();
        aux = aux. Desapila();
        exp. Apila( eval(x, op, y) );
    } else { aux. Apila( exp. Tope() );
        exp = exp. Desapila();
    }
} catch (...) { ... }
}

```

Huffman 2023 | coded, receives Huffman tree "t" and ASCII code of symbol "s", returns lists of 0 and 1 the Huffman code of "s", if included. If not, [-1]

```
public static Lista<Integer> code01(Arbol Binario<Integer> t, Integer s)
{
    Lista<Integer> l, laux; l = new LD();
    Arbol Binario<Integer> aux = new ABD();
    try { if (t.subArbolDer().EsVacio())
        if (t.subArbolIzg().Raiz() != s) l.Añade(-1);
    } else { l = code01(t.subArbolIzg(), s);
        if (l.EsVacio() || l.Cabeza() != -1) l.Añade(0);
        else { l = code01(t.subArbolDer(), s);
            if (l.EsVacio() || l.Cabeza() != -1) l.Añade(1);
        }
    }
    { catch (-) } ...
    { return l;
    }
}
```

tdv
tir s -1

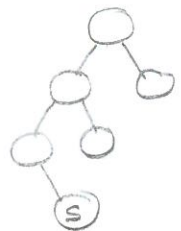
Huffman 2022 | percent. receives Huffman tree and a list of 0s and 1s. If the list corresponds to the Huffman code of a symbol in the tree, return %, else, -1

```
public static int percent(Arbol Binario<Integer> t, lista<Integer> l)
{
    integer p = -1; AB<Integer> aux = new ABD();
    if (!t.EsVacio())
    {
        "final" } try { if (l.EsVacio()) {
            if (t.subArbolDer().EsVacio()) p = t.Raiz();
            else if (l.cabeza() == 0)
                p = percent(t.subArbolIzg(), l.cola());
            else { p = percent(t.subArbolDer(), l.cola());
            }
        } catch (TADVacio Exception ex) { "ex" }
    }
    { return p;
    }
}
```

Hemos llegado al final del listado, ese era el símbolo

se devuelve "s"

0 0 1



1. Thread 1 ⇒ L no vacía

2. Primero un 0, se crea Thread 2 (i)

Thread 2 ⇒ L no vacío

Segundo un 0, se crea Thread 3 (i)

Thread 3 ⇒ L no vacío

Tercero un 1, se crea Thread 4 (d)

Thread 4 ⇒ L es vacío

No hay hijo derecho → El nodo es el buscado

p = t.raiz()

SOLUTIONS

1)

```
Pila s = new PD(); s.APila(3); s.APila(5); s.APila(12); s.APila('%');
s.APila('*'); s.APila(2); s.APila('+');
```

ITE. 1	ITE. 2	ITE. 3	ITE. 4	ITE. 5	ITE. 6	ITE. 7	ITE. 8	ITE. 9	ITE. 10
2									
'*'		'*'			12				
'%'		'%'		'%'	'%'		2		
12		12		12	'*'	12	'*'		'*'
5		5	2	5	2	5	2	2	2
3	'+'	3	'+'	3	'+'	3	'+'	3	'+'
exp	aux	exp	aux	exp	aux	exp	aux	exp	aux

2)

```
public static int percent ( ArbolBinario<Integer> t,
                           Lista<Integer> l)
{
    Integer p=-1; ArbolBinario<Integer> aux=new ABD<>();
    if (!t.EsVacio())
        try{ if (l.EsVacia())
                {if (t.SubArbolDcho().EsVacio()) p=t.Raiz();}
                else if (l.Cabeza()==0)
                    p=percent(t.SubArbolIzqdo(),l.Cola());
                else p=percent(t.SubArbolDcho(),l.Cola());
            }catch(TADVacioException e) {System.out.println(e);}
    return p;
}
```

3)

```
public static Lista<Integer> cycle(int[][] adj) {
    int inf=Integer.MAX_VALUE, minweight=Integer.MAX_VALUE,
    n= adj.length, i=0, j=0;
    Lista<Integer> sol=new LD<Integer>();
    Lista<Integer>[][] path = Floyd(adj);
    for (i=0; i<n; i++)
        for (j=0; j<n; j++)
            if (i!=j && adj[i][j]!=inf && adj[j][i]!=inf
                && (adj[i][j]+adj[j][i])<minweight)
            { minweight=adj[i][j]+adj[j][i];
              sol=(Lista<Integer>)(path[i][j].Cola().clone());
              sol.Concatena(path[j][i].Cola()); }
    return sol;
}
```

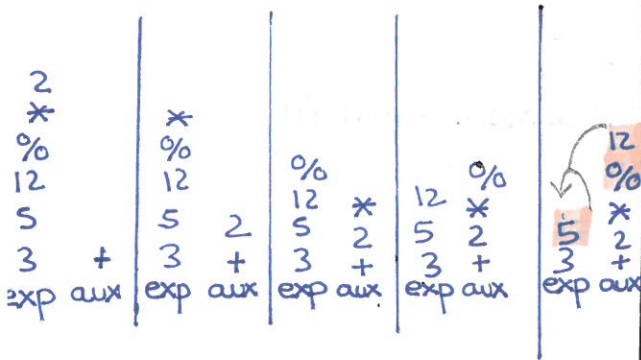

Polish Notation

$$2 + ((12 \% 5) * 3)$$

PN ←

2 + ((12 % 5) * 3)
 + 2 ((12 % 5) * 3)
 + 2 (x (12 % 5) 3)
 + 2 (x (% 12 5) 3)
 + 2 x % 12 5 3

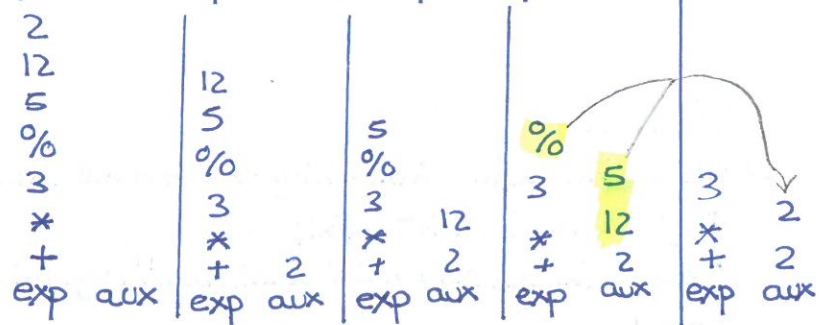
Se apila de derecha a izquierda
 s.apila(3); s.apila(5); s.apila(12);
 s.apila(%); s.apila(*); s.apila(2);
 s.apila(+);



↻ RPN (Reverse)

2 + ((12 % 5) * 3)
 2 ((12 % 5) * 3) +
 2 ((12 % 5) 3) * +
 2 ((12 5) % 3) * +
 2 12 5 % 3 * +

Se apila de derecha a izquierda
 s.apila(+); s.apila(*); s.apila(3); s.apila(%);
 s.apila(5); s.apila(12); s.apila(2);



Se va desapilando en aux, cuando en aux hay dos n° seguidos y un operador en exp, se opera y se guarda en aux

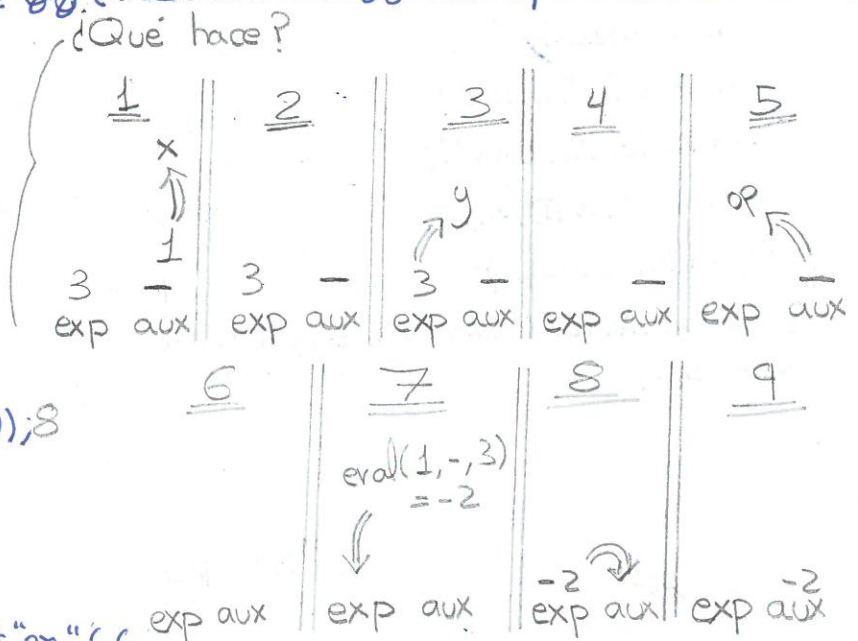
Código

Pila exp ⇒ tiene la expresión en PN, ex: + 2 * 3 4 ⇒ 2 + (3 * 4)

Double ⇒ guarda números; String ⇒ guarda operadores

Proceso: Se recorre exp de tope a abajo, evaluando posibilidades. Aux: pila que ayuda

```
public static Double PNvalue(Pila exp) {
    Pila aux = new PD();
    Double x, y;
    String op;
    try {
        while (!exp.EstVacia()) {
            if (exp.Tope() instanceof Double && (!aux.EstVacia() && aux.Tope() instanceof Double)) {
                // ¿Qué hace?
                x = (Double) aux.Tope();
                aux = aux.Desapila();
                y = (Double) exp.Tope();
                exp = exp.Desapila();
                op = (String) aux.Tope();
                aux = aux.Desapila();
                exp.Apila(eval(x, op, y));
            } else {
                aux.Apila(exp.Tope());
                exp = exp.Desapila();
            }
        }
        return (Double) aux.Tope();
    } catch (TADVacioException) {
        sout "ex" &&
    }
}
```

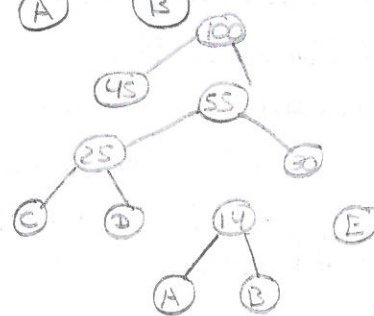
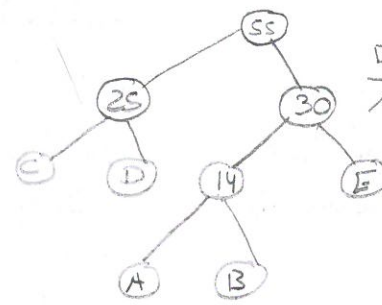
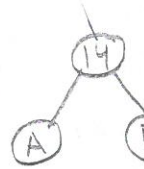
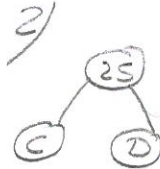
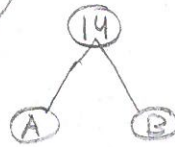
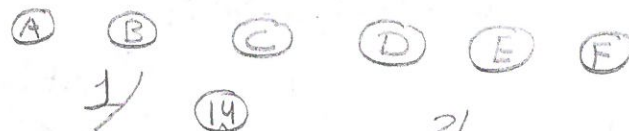


Huffman Trees

De menor a mayor

Símbolo	Freq
A	5
B	9
C	12
D	13
E	16
F	45

Tree Creation



Método insert

```

public static Lista<ABD<Integer>> insert(ABD<Integer> x, Lista<ABD<Integer>> y)
{
    if (y. EsVacia()) { y. Añade(x); }
    else if (x. Raiz() <= y. Cabeza(). Raiz()) { y. Añade(x); }
    else
    {
        ABD<Integer> c = y. Cabeza();
        y = y. Cola();
        y = insert(x, y);
        y. Añade(c);
    }
    return y;
}

```



Método Huffman Tree

```

public static <ABD<Integer>> HuffmanTree(Lista<ABD<Integer>> x)
{
    ABD<Integer> sol = new ABD<Integer>;
    try { while (!x. Cola(). EsVacia()) {
        ABD = new ABD<Integer>(x. Cabeza(). Raiz() + x. Cola(). Cabeza(). Raiz()); // suma de Cabezas
        x. Cabeza(). Raiz() -> El 1º elem. Hijo izq
        x. Cola(). Cabeza(). Raiz() -> El 2º elem. Hijo der.
        x = x. Cola(). Cola(); -> Pasa de los dos siguientes ya que fueron procesados
        x = insert(ABD, x); -> Va a insert para asegurar orden creciente
        sol = x. Cabeza();
    }
    } catch (TADvacuoException ex) { xat(ex); }
    return sol;
}

```

Lista **ORDENADA** creciente

NAME: _____

Java interfaces that can be used in this exam

```
public interface Lista<E> {
    public boolean EsVacia();
    public void Añade(E x);
    public E Cabeza() throws TADVacioException;
    public Lista<E> Cola() throws TADVacioException;
    public void Concatena(Lista<E> l);
    public Object clone();
}

public interface ArbolBinario<E> extends Arbol<E>
{
    public boolean EsVacio();
    public E Raiz() throws TADVacioException;
    public ArbolBinario<E> SubArbolIzqdo() throws TADVacioException;
    public ArbolBinario<E> SubArbolDcho() throws TADVacioException;
}
```

(12%5)*3
↓
*(12%5)3
↓
*%1253

1) [Polish notation] In this exercise we are concerned with the second practice devoted to evaluate an arithmetic expression stored into a general stack once expressed in PN (*Polish notation*). Given the arithmetic expression $2 + ((12 \% 5) * 3)$, complete the following Java code for initializing a stack representing the same expression in Polish Notation: $+ 2 * \% 12 5 3$

```
Pila s = new PD(); s.APila( 3 ); s.APila( 5 ); s.APila( 12 );
s.APila( % ); s.APila( * ); s.APila( 2 ); s.APila( + );
```

In the table below, use only the cells required to represent the contents of stacks **exp** and **aux** after finishing each iteration when calling the method explained in class with the previous stack:

ITE. 1	ITE. 2	ITE. 3	ITE. 4	ITE. 5	ITE. 6	ITE. 7	ITE. 8	ITE. 9	ITE. 10
2									
*	*			12					
%	%	%	%	%		2			
12	12	12	*	12	*	*	*		
5	5	2	5	2	5	2	2	2	
3	+	3	+	3	+	3	+	6	8
exp	aux	exp	aux	exp	aux	exp	aux	exp	aux

2) [Huffman trees] Complete in the next page the Java code of method **percent** which receives a Huffman tree with at least two leaves and a list only containing 0's and 1's. If such list corresponds to the Huffman code of a symbol in the tree, then the method must return the percentage of such symbol. Otherwise, the output will be -1.

3) [Floyd's algorithm] A cycle in a graph is a path starting and ending in the same node and the rest of nodes are not repeated. Complete in the next page the Java code of method **cycle** which receives the adjacency matrix of a graph and, after applying the Floyd's algorithm on it, returns a list with all the nodes (without repetitions) involved on the cycle with the lowest weight in the graph. If there are no cycles in the graph, the output will be an empty list.

```
public static int percent (ArbolBinario<Integer> t, Lista<Integer> l)
{   Integer p=-1; ArbolBinario<Integer> aux=new ABD<>();
    if (!t.EsVacio())
        try{
```

```
        }catch(TADVacioException e) {System.out.println(e);}
    return p;
}
```

```
public static Lista<Integer> cycle(int[][] adj)
{   int inf=Integer.MAX_VALUE, minweight=Integer.MAX_VALUE;
    n= adj.length, i=0, j=0, k=0;
    Lista<Integer> sol=new LD<Integer>();
    Lista<Integer>[] path = Floyd(adj);
```

```
    return sol;
}
```

SOLUTIONS

1)

```
Pila s = new PD(); s.APila("+"); s.APila("*"); s.APila(3.0);
s.APila("%"); s.APila(5.0); s.APila(12.0); s.APila(2.0);
```

ITE. 1	ITE. 2	ITE. 3	ITE. 4	ITE. 5	ITE. 6	ITE. 7	ITE. 8	ITE. 9	ITE. 10
12.0									
5.0	5.0								
“%”	“%”	“%”							
3.0	3.0	3.0	5.0	3.0					
“*”	“*”	12.0	“*”	12.0	“*”	2.0	“*”	2.0	6.0
“+”	2.0	“+”	2.0	“+”	2.0	“+”	2.0	“+”	2.0
exp	aux	exp	aux	exp	aux	exp	aux	exp	aux

2)

```
public static Lista<Integer> code01(ArbolBinario<Integer> t, Integer s)
{
    Lista<Integer> l, laux; l=new LD();
    ArbolBinario<Integer> aux = new ABD<>();
    try {
        if (t.SubArbolDcho().EsVacio())
            { if (s != t.SubArbolIzqdo().Raiz()) l.Añade(-1); }
        else { l = code01(t.SubArbolIzqdo(), s);
            if ((l.EsVacia() || l.Cabeza() != -1)) l.Añade(0);
            else { l = code01(t.SubArbolDcho(), s);
                if ((l.EsVacia() || l.Cabeza() != -1)) l.Añade(1);
            }
        }
    } catch (TADVacioException e) { System.out.println(e); }
    return l;
}
```

3)

```
public static int flodijs(int[][] adj) {
    int inf = Integer.MAX_VALUE, n = adj.length, i = 0, j = 0, k = 0, num = 0;
    int[][] paw = (int[][]) adj.clone(); Lista<Integer>[] pan = Floyd(paw);
    for (i = 0; i < n; i++)
        for (j = 0; j < n; j++)
            try {
                if (i != j && adj[i][j] == inf && paw[i][j] != inf &&
                    (2 * adj[i][pan[i][j].Cola().Cabeza()]) < paw[i][j]) num++;
            } catch (TADVacioException e) { System.out.println(e); }
    return num;
}
```


if $[(H_i, leaf_i) > (H_i, leaf_d)] = H_i, leaf_i$
 if $[(H_i, leaf_i) < (H_i, leaf_d)] = H_i, leaf_d$

$dec24(AB \perp (ABS \vee (ABS \vee AV)) \vee) \lfloor 7 \rfloor =$

$= dec24 i (dec24 d \perp) = dec24 i (dec24 AV \lfloor 7 \rfloor) = dec24 i \lfloor 7 \rfloor$

$= dec24 (ABS \vee (ABS \vee AV)) \lfloor 7 \rfloor = dec24 i (dec24 d \perp) =$

$= dec24 i (dec24 (ABS \vee AV) \lfloor 7 \rfloor) = dec24 i (8:7) = dec24 i \lfloor 8:7 \rfloor$

$= dec24 AV \lfloor 8:7 \rfloor = \lfloor 8:7 \rfloor$

if (!t.EsVacio())

try { t1.subord b9.EsVacio();
i++;

fnodes(t1.subord b9())

sol--;

if (-sol % (2i+1) == 0)

sol = Math.abs(sol); {

sol



$2i+1$

2) 6

Node EOJ: new ^{Node} Gic (o(x, Rotz(i), (d x)

[numnodes - 1] [numnodes]

Graphs: High Level

public static <E> boolean IsConnected (Grafo <E> g)

Lista <E> nodes, visited, pending; nodes = g.Nodos();
integer toVisit = nodes.Longitud(); if toVisit == 0 return true

try { pending.Add(nodes.Cabeza()); // 1º elem. de la lista

while (!pending.EsVacia() || toVisit > 0) // Solo para la 1ª iteración

E x = pending.Cabeza(); pending = pending.Cola();

if (!visited.EstaContenida(x)) { visited.Add(x); toVisit--; pending.Concatena(g.Adya-centes(x));

x.visitado = true

Lista <E> aux = g.Adya-cente(x)

aux.Concatena(pending)

pending = aux

GNDE => Solo este tiene el visitado Atributo

{ catch (...) }

return toVisit == 0;

High ++ Low Level

public static <E> boolean IsConnected (GNDE <E> g)

Lista <E> Nodes, pending; nodes[] = g.Nodos;
int toVisit = nodes.Longitud

if (x.visitado

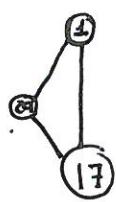
try { pending.Add(nodes.Cabeza());

while (!pending.EsVacia() || toVisit > 0)

E x = pending.Cabeza(); pending = pending.Cola();

if (!visited.EstaContenida(x)) { visited.Add(x); toVisit--; pending.Concatena(g.Adya-centes(x));

Graphs



Limited size:

MAX-nodes = 5

num Nodes = 3

	0	1	2	3	4
0		T	T		
1	T		T		
2	T	T			
3					
4					

↓ ↓ ↓ (Nodes Name)
1 29 17
↓ ↓ ↓ (Visited?)
F F F

En BST el mínimo está en 



```

graph LR
    Mixheap --> Mayor
    Mixheap --> Menor

```

Java Iterative inside GND class

public GNDE (Arbol Binario <integer> t) {
 Diamonds = 0;
 ...

```
numnodos = 0; Adyacencias = new boolean [Max_Nodos] [Max_Nodos];
Nodos = new Node [Max_Nodos];
```

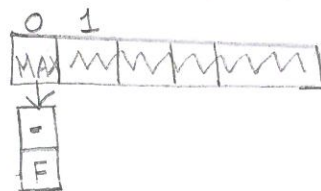
Nodos = new NodoGrafo [Max_Nodos]; Integer x; Arbol Binario <Integer> y, z;

```
try { if (!t, EsVacio()) }
```

Nodos[0] = NovoGrafo(t.Raiz(), false);

```
t = mixheap(t, subA[2q], t, subA[2r-1]);
```

flum nodos + r;

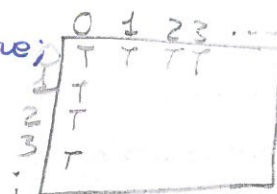
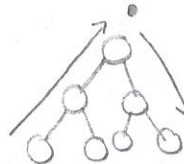
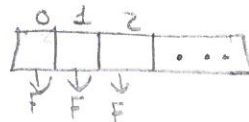


```
while (!EsVacio())
```

Nodos [numnodos]: Nodo Grafo (t. Raiz (), false);

```
t = mixheap(t.subAlg(), t.subADer());
```

Adyacencia [0] [numnodos] = Adyacencia [numnodos] [0] = true

$$\text{num nodes} + 1$$


```
{ catch (catch TADVacuoException ex) sout(ex); }
```

public GND E (Arbol Binario $\langle E \rangle^+$).

Adyacencia = New Boolean [Max-Node] [Max-Node]; Nodes = new Nodo Grafo [Max-Node];

```
if (!t.EsVacio())
```

```
try { Nodos[0] = new Nodo Grafo(t.Raiz(), false);
```

modografo τ ; auxz t. subarb $129(1)$;

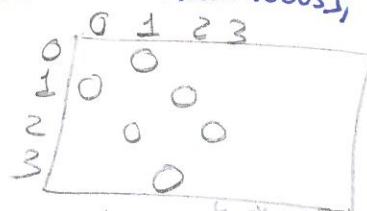
```
while (!aux. EsVacio()) {
```

Nodos [nuevnodos] = new Nodografo (aux. Raiz (), false);

Adjacencia [numeros - 1] [numeros] : Adj [numeros] [numeros - 1] = true;

newNode; aux = t.subgraph(1);

if (!subArb12 q(). EsVao0()))



→ Esto corresponde
el con el anterior
nodos - 1] = true;

Unit 4. Trees

Árbol: Colección de nodos conectados jerárquicamente

Raíz: Nodo principal del árbol

Hijo: Nodos conectados a otro previo

Hojas: Nodos sin hijos

Subárbol: Cualquier nodo + sus descendientes

Nivel: Profundidad de un nodo, raíz en nivel 1 (o 0)

Altura: Nivel más profundo del árbol

Recorrido: Proceso de visitar todos los nodos

- Preorden: Raíz, subárbol izq, subárbol der.
- Inorden: Subárbol izq, raíz, subárbol der.
- Postorden: Subárbol izq, subárbol der, raíz.

empty :: Tree a → Bool

root :: Tree a → a

leftsubtree // rightsubtree :: Tree a → Tree a

Implementar recorridos

inorder :: Tree a → [a]

inorder EBT = []

inorder (BT r i d) = inorder i ++ [r] ++ inorder d

preorder :: Tree a → [a]

preorder EBT = []

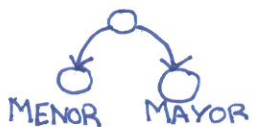
preorder (BT r i d) = [r] ++ preorder i ++ preorder d

postorder :: Tree a → [a]

postorder EBT = []

postorder (BT r i d) = postorder i ++ postorder d ++ [r]

Árboles de búsqueda (Binary search trees)



→ Queremos introducir X
 → if $X < r$ = izq.
 → if $X > r$ = der.

insert x EBT = BT x EBT EBT
 insert x (BT r i d)
 | $x < r$ BT r (insert x i) d
 | $x > r$ BT r i (insert x d)

Search x EBT = False

Search x (BT r i d)

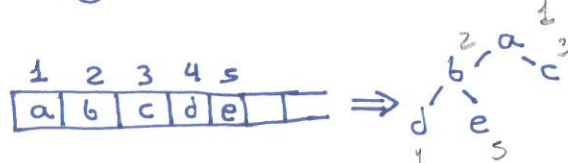
| $x == r$ ⇒ True

| $x < r$ = BT r (search i) d

| $x > r$ = BT r i (search d)

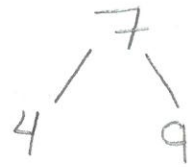
height EBT = 0

height (BT r i d) = 1 + max (height i) (height d)



$2i \leftarrow 0 \rightarrow 2i+1$

$2i+1 \leftarrow 1 \rightarrow 2i+2$



mytree = BT 7 (BT 4 EBT EBT) (BT 9 EBT)

↳ mytree = BT 7

(BT 4

EBT

EBT)

(BT 9

EBT

EBT)

Vacio

data Tree a = EBT | BT a (Tree a) (Tree a)

1-Cont. long Esta conectado?

Nodo con raíz a subárbol izq y der.

Meto nodos, nodes = g.Nodes();

Tomo cabeza, peroing = nodes.Cabeza();

Miro si recorren

AVL Height AVL

HeightAVL :: Tree a $\rightarrow$ int

HeightAVL EBT = 0

HeightAVL (BT_ i d) = left = HeightAVL i ; right = HeightAVL d ;

if left == -1 || right == -1 || abs(left - right) > 0

then -1

else 1 + max(left) (right)

InvertadoBT :: Tree a $\rightarrow$ [a]

El árbol es t

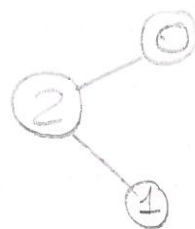
InvertadoBT EBT = [] // No hay na

InvertadoBT (BT_ i d) = left = InvBT i ; right = InvBT d ;

if left - right != 0 = 0 // No es completo

else r: (BT + i) : (BT + d)

0 0 1 2 3
1
2 T T T
3



Practica de Floyd

Implement public static Lista <Integer>[][] Floyd(int[][] adj)

Shortest path in a directed graph

TAD COLA

cabeza → 4, 3, 2, 3, 32, 0, 4, 7 → último elemento introducido
cola o queue

Sigue la idea de FIFO, first In, first Out
La cabeza fue el primero metido y el primero en salir

data Queue a = EQIPut (Queue a) a

Free generators

EQ :: 2 :: 7 :: 5 :: 7

Put == :

EQ: vacía cola
empty queue

Operaciones

• EmptyQ :: Queue a → Bool

EmptyQ EQ = true

• first :: Queue a → a

① first (EQ :: x) = x

② first (c :: x) = first c

EXAMPLE

first (EQ :: 3 :: 5 :: 2 :: 7)

↳ 3

② → 2

first (EQ :: 3 :: 5 :: 2 :: ...) ; first (EQ :: 3)

② → 5

① → 3

• rest :: Queue a → Queue a

① rest (EQ :: x) = EQ

② rest (c :: x) = (rest c) :: x

EXAMPLE

rest (EQ :: 3 :: 5 :: 2 :: 7)

↳ EQ :: 5 :: 2 :: 7

TAD STACK (O pila)

3 → Tope o cima

2
5
2
7 } Cola, desapila o pop

Para poder acceder a los elementos de abajo se debe desapilar

ES: PV: pila vacía

data Pila a = PV | a :/ (Pila a)

:/ apila o push
PV pila vacía

3:/2:/5:/2:/7:/ES

• top :: Stack a → a Like head

• pop :: Stack a → Stack a Like tail

⚠ Típico error ⚠

elimina :: a → Stack a → Stack a

elimina x ES = ES

elimina x (y:/p) if x == y then p else [y:/ (elimina x p)] → Elimina de abajo sin desapilar

Podemos añadir

"apilar :: a → Stack a → Stack a Quedando
apilar x p = x:/ "

elimina :: a → Stack a → Stack a

elimina x ES = ES

elimina x (y:/p) if x == y then p else apilar y (elimina x p)

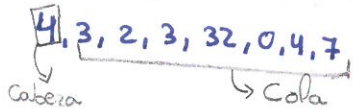
Tema 3. Listas. Pilas. Colas

✦ Tipo abstracto de datos (TADs):

→ Lineal: vacío o al menos un elemento siguiendo un orden
↳ Arrays, listas, pilas, colas

→ No lineal: conjuntos, árboles, grafos

TAD LISTA

 La cabeza es último elemento introducido
Sigue una estructura LIFO (Last IN First Out)

Todas son de un mismo tipo pero no de tipo fijo → Genérica

Si se mezclan tipos → general

Definiciones Haskell

data Lista a = Nil | cons a (Lista a)

Nil: lista vacía

Cons: operaciones generadoras libres

Lista a = [a]

Nil = []

Cons " = (infijo)

Operaciones no generadoras

Cons 1 (Cons 2 Nil) = 1 : (2 : [])
['a', 'b', 'c'] = "abc"

empty :: [a] → Bool

empty [] → false

empty _ = False

head :: [a] → a

head [] → error "Empty list"

head (h : _) = h

tail :: [a] → a

tail [] → error "Empty list"

tail (_ : t) = t

length :: [a] → Int

length [] = 0

length (h : t) = 1 + length of t

concatenate :: [a] → [a] → [a]

↳ also seen as ++ [a]
"Link two lists"

concatenate [] s = s

concatenate (h : t) s = (h : concatenate t s)

reverse :: [a] → [a]

↳ Same elements but in inverted order
↳ Example: reverse [1,2,3] } h = 1
t = [2,3]

reverse [] = []

reverse (h : t) concatenate (reverse t) [h]

reverse [3] is [] ++ [3] = [3]; reverse [2,3] is [3] ++ [2] = [3,2];
reverse [1,2,3] is [3,2] ++ [1] = [3,2,1]

last :: [a] → a

"Returns last element of a list"

last [x] = x

([x] == to x: [], list with 1 element)

last (_ : t) = last t

elem :: (Eq a) ⇒ a → [a] → Bool

"If element in list" "a belongs to Eq class to compare"

extract :: (Eq a) ⇒ a → [a] → [a]

extract x [] = []

extract x (h : t) = if (x == h) then extract h, else h : extract x t

take :: Int → [a] → [a]

take 0 [] = []

take n (h : t) = h : take (n-1) t

→ ya que h se "salva"

cut :: Int → [a] → [a]

cut 0 1 = 1

cut n (h : t) = cut (n-1) t + h se fue

substitute :: (Eq a) ⇒ a → a → [a] → [a]

↳ "changes all x in l by y"

substitute x y (h : t) if (h == x) then y substitute x y t / else h : substitute x y t

sum :: [Int] → Int

sum [] = 0

sum (h : t) = h + sum t

greater :: (Ord a) ⇒ [a] → a

greater [x] = x

greater (x : y : t) = if x > y then greater (x : t) else greater (y : t)

23-24

1st Progress test

op23 :: (ord a) => [a] -> [a] Doy una lista ordenada y recibo una lista

op23 (x:y:s) = if (x > y) then s ++ [x] else x: (op23 (y:s))

op23 1 = 1

a) op23 [1,2,3] -> [1,2,3]

b) op23 [3,2,1] -> [1,3]

(3 > 2) -> 1 ++ [3] => [1,3]

op23 1 = 1 -> 1

c) op23 [1,3,2] -> [1,3]

(3 > 2) -> [] ++ [3] = [3]

try}... { catch (TADVacioException ex) { ex.printStack Trace(); {

In Lista <E>:

public boolean EsVacia();

public void Añade(E e);

public E Cabeza() throws TADVacioException;

public Lista<E> Cola() throws TADVacioException;

In LD<E>:

private int longitud;

private Nodo<E> NodoCabeza

In Nodo <E>:

public Nodo (E e, Nodo <E> n) { Info = e; Siguiente = n; {

a) public LD (Lista <E> L1, LD <E> L2) {

Nodo <E> aux = null

public LD (Lista <E> L1, LD <E> L2) {

longitud = L2.longitud; NodoCabeza = L2.NodoCabeza;

try { while (!L1.esVacia()) {

NodoCabeza = new Nodo (L1.Cabeza (), NodoCabeza);

L1 = L1.Cola(); longitud++;

{ catch (TADVacioException ex) { ex.printStack Trace(); { {

b) public static <E extends Comparable> Lista <E> op23 (Lista <E> L)

{ E x, y; LD <E> aux1, aux2; aux1 = new LD <E> (); aux2 = new LD <E> ();

if (L.esVacia() || (L.Cola().esVacia())) return L;

try { x = L.Cabeza(); l = L.Cola(); y = l.Cabeza();

while (!l.Cola().esVacia() || x.compareTo(y) < 0) {

aux1.añade(x); x = y; y = l.Cabeza(); l = l.Cola(); {

if (x.compareTo(l.Cabeza()) > 0) { aux2.añade(x); l = l.Cola(); {

else aux1.añade(x);

while (!l.esVacia()) {

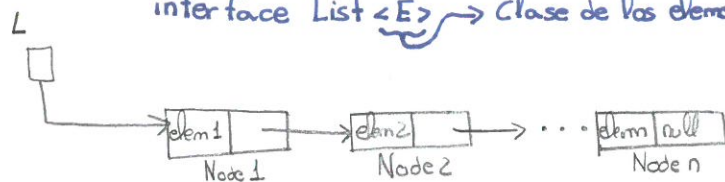
aux1.añade(l.Cabeza()); l = l.Cola(); {

{ catch (TADVacioException e) { return new LD <E> (); {

return new LD <E> (aux1, aux2); {

Implementación de listas

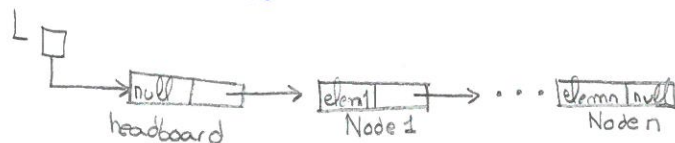
interface List $\langle E \rangle$ $\rightarrow$ clase de los elementos de la lista



Nodes from Node $\langle E \rangle$ class

L instance of List $\langle E \rangle$ interface
elems to be stored

$\rightarrow$ With headboard



Para evitar asimetría con el primer nodo, ya que sería el único sin predecesor y obligaría a modificar algoritmos

Insert::

Void Insert($E x, E y$) throws TADVacioException. Primero y , luego x

insert:: $a \rightarrow a \rightarrow \text{Stack } a \rightarrow \text{Stack } a$
(Eg a)

insert $y x$ EmptyS = ES

insert $y x$: if $x == a$ then Node y (Node a rest) else Node a (insert $x y$)

TEORIA

80%

Tema }
Listas, Pila, Cola }

1
0
7

• Tema 4
Arboles

1
0
7

TEMAS
Grafos.

50% JAVA.

30% HASKELL

PRACTICAS

20%

JAVA texto o video...
Grafos.

107 ejercicios o preguntas

20% JAVA.

NAI Polare — PILAS
Huffman — ARBOLES
Floyd. — GRAFOS

Updated in line with Cambridge English: First (FCE) 2015 revisions



Roy Norris
3rd Edition
coursebook with key

Ready for First

MACMILLAN EXAMS 

```
public void Inserta (E x, E y) throws TADVacioException {
```

```
    Nodo <E> aux;
```

```
    boolean encontrado = false;
```

```
    if (EsVacio()) throw new TADVacioException
```

```
    if (Nodo Cabeza. Info. equals(x)) {
```

```
        Nodo Cabeza = new Nodo <E> (y, Nodo Cabeza);
```

```
        longitud++;
```

```
    } else {
```

```
        aux = Nodo Cabeza;
```

```
        while ((aux. Siguiente != null) && ! encontrado)
```

```
            if (aux. Siguiente. Info. equals(x)) {
```

```
                aux. Siguiente = new Nodo (y, aux. Siguiente);
```

```
                longitud++;
```

```
                encontrado = true;
```

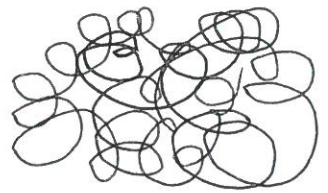
```
            } else { aux = aux. Siguiente;
```

```
            }
```

```
    }
```


Estr. lineales Java

↳ árboles al final



/

Grados Haskell → No

Proyecto Java Temos 3, 4, 5
Prácticos

Árboles → Nada de generadores, solo binarios y sus
variaciones

Números no mucho en los exámenes

Bajo nivel no llamo otros funciones

Alto nivel no pointers

Parcial 1 → 3

Parcial 2 → 4, 5 1a 22-23 No, que llea grafos

13 Animal magic

Vocabulary 1: The Arts

- 1 Both words in each of the pairs below can be used in combination with one of the words in the box. Write an appropriate word from the box in each of the spaces. There is an example at the beginning (0).

novel	opera	concert	painting	stone	classical	gallery
-------	-------	---------	----------	-------	-----------	---------

- | | | | | | |
|--------------------|----------------|---------|---------------------|---------------------------|-------|
| 0 open-air
jazz | <u>concert</u> | 1 _____ | ballet
music | 4 portrait
art | _____ |
| | | 2 _____ | singer
house | 5 abstract
priceless | _____ |
| | | 3 _____ | sculpture
statue | 6 detective
historical | _____ |

- 2 Which people do you associate with each of the following areas of the arts?

theatre	music	literature	art	opera	ballet	sculpture
---------	-------	------------	-----	-------	--------	-----------

Example:

Theatre: actor, actress, director, cast, playwright, audience

Check your answers in the Wordlist on page 208.

Reading and Use of English

Part 6

Gapped text

- 1  Look at the work of art in the photograph and discuss the following questions.
What type of person might want to own such a work of art?
Where might they put it on display?

